Teaching Rigorous Distributed Systems With Efficient Model Checking

Ellis Michael

Doug Woos

Thomas Anderson

Michael D. Ernst

Zachary Tatlock

UNIVERSITY of WASHINGTON











UW CSE 452

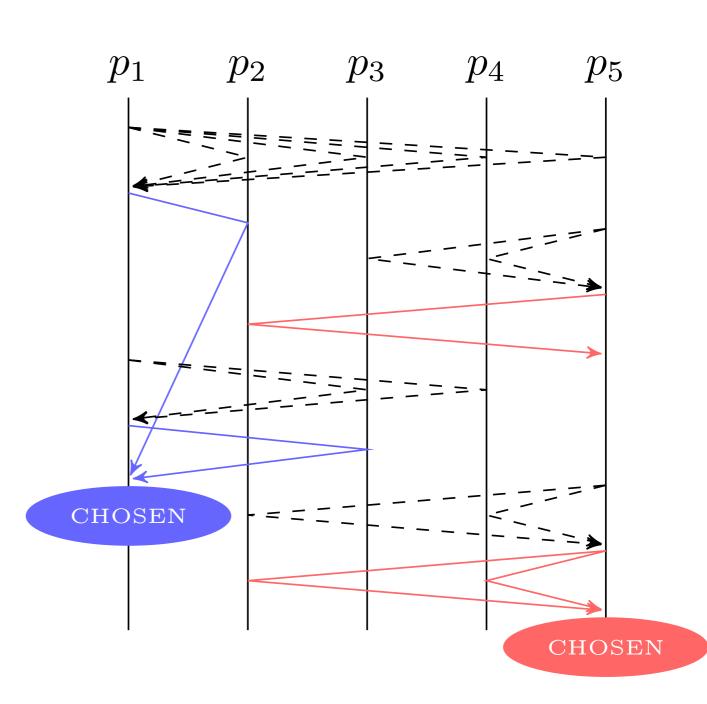
- Course on distributed systems for undergraduates and 5th year
 Master's students, enrollment grown to approximately 200
- Lab assignments building fault-tolerant, consistent distributed systems, based on assignments developed for MIT 6.824:
 - 1. Exactly-once RPC
 - 2. Primary-backup
 - 3. Paxos-based state machine replication
 - 4. Sharded key-value store
 - 5. Distributed transactions using two-phase commit
- Tests used for grading assignments given to students

Goal: Tests which identify common bugs, provide timely feedback, and assist debugging to help students build systems to rigorous standards.

Systems solution for teaching distributed systems

Testing Distributed Systems is Difficult

- Simple Paxos bug: leader checks for quorum with matching values (rather than proposal numbers).
- Finding such a bug is difficult with current tools.
- This false quorum bug could be caused by a fundamental misunderstanding.



"Just 3 days before the deadline of the project, my partner and I discovered that our Paxos failed 1 of 100,000 tests. ...We realized that the bug comes from our optimization of duplicate request detection before putting request on the Paxos operation log. ... We needed to rewrite fifty percent of the whole project but we did not give up. Finally, after 30 hours of work in 2 days, we fixed the design flaw and eliminated the bug. We were so excited that we started to dance in the lab."

- CSE 452 Student

Checking Correctness

- Execution-based testing is insufficient; can miss bugs unlikely to occur based on timing.
- Manual review does not scale or provide feedback quickly enough.
- Formal verification is difficult and time-consuming, not approachable for students.

Checking Correctness: Model Checking

- Researchers and practitioners use model checking to validate protocols and software, systematically searching through possible executions.
- Some specification languages are difficult to learn, do not produce runnable code.
- Naïve methods do not scale well, fail to find rare bugs quickly and reliably.

DSLabs

A framework for creating distributed systems labs and test suites

... capable of finding common bugs in students' implementations quickly and reliably

... using a widely-used programming language (Java) and easily-learned tools

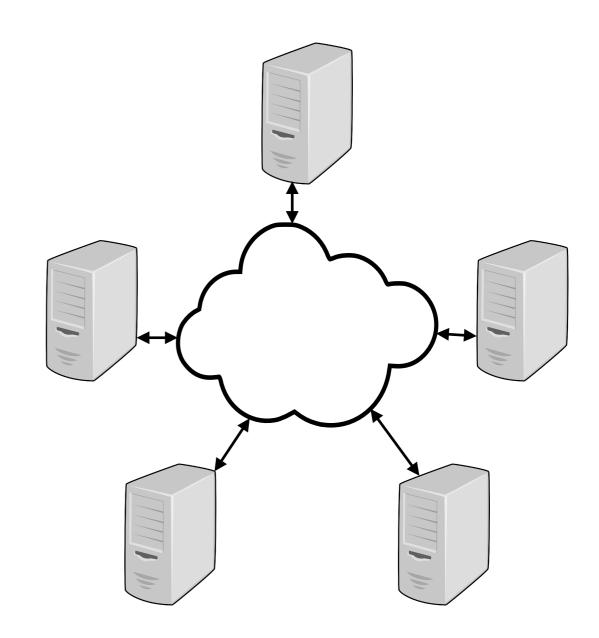
... that helps students write correct, efficient, runnable code

... and understand errors when they do arise.

The Rest of This Talk

- 1. The DSLabs programming model
- 2. Model checking strategies and optimizations
- 3. Understandability and Oddity visual debugger
- 4. Experiences

- A distributed system consists of a set of nodes which communicate over an asynchronous network, working together to run a protocol.
- Nodes are I/O automata; they run as single-threaded event loops.
- Nodes are split between client and server nodes.



- A distributed system consists of a set of nodes which communicate over an asynchronous network, working together to run a protocol.
- Nodes are I/O automata; they run as single-threaded event loops.
- Nodes are split between client and server nodes.



```
1: init()
```

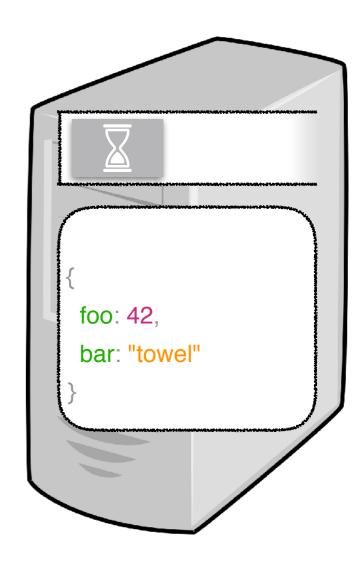
2: loop

4: update_state(e)

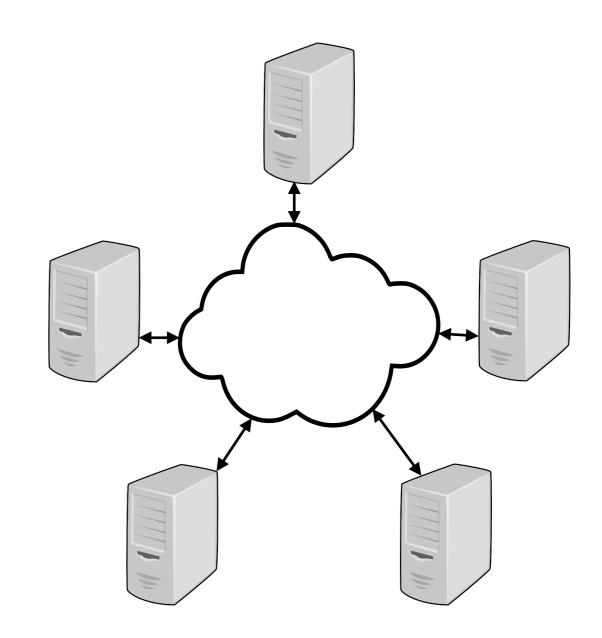
5: send_msgs()

6: set_timers()

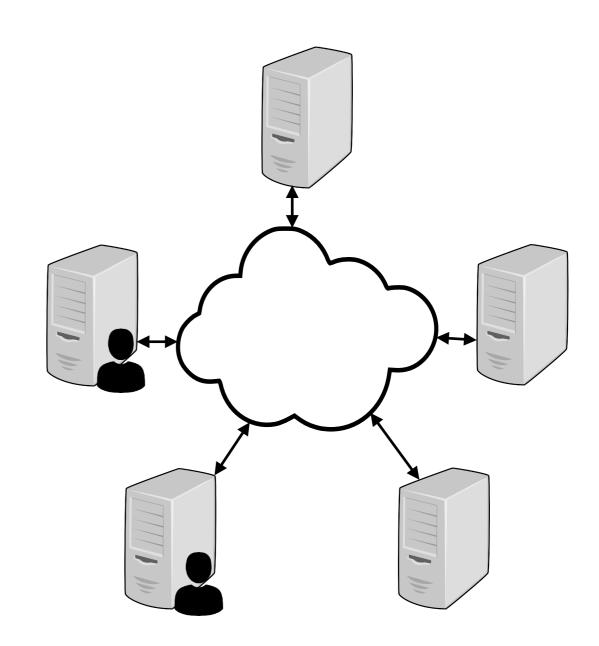
7: endloop



- A distributed system consists of a set of nodes which communicate over an asynchronous network, working together to run a protocol.
- Nodes are I/O automata; they run as single-threaded event loops.
- Nodes are split between client and server nodes.



- A distributed system consists of a set of nodes which communicate over an asynchronous network, working together to run a protocol.
- Nodes are I/O automata; they run as single-threaded event loops.
- Nodes are split between client and server nodes.

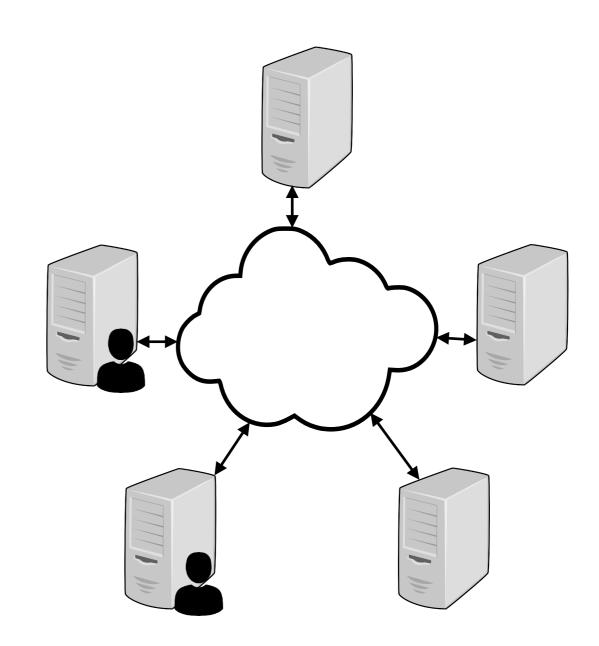


- A distributed system consists of a set of nodes which communicate over an asynchronous network, working together to run a protocol.
- Nodes are I/O automata; they run as single-threaded event loops.
- Nodes are split between client and server nodes.



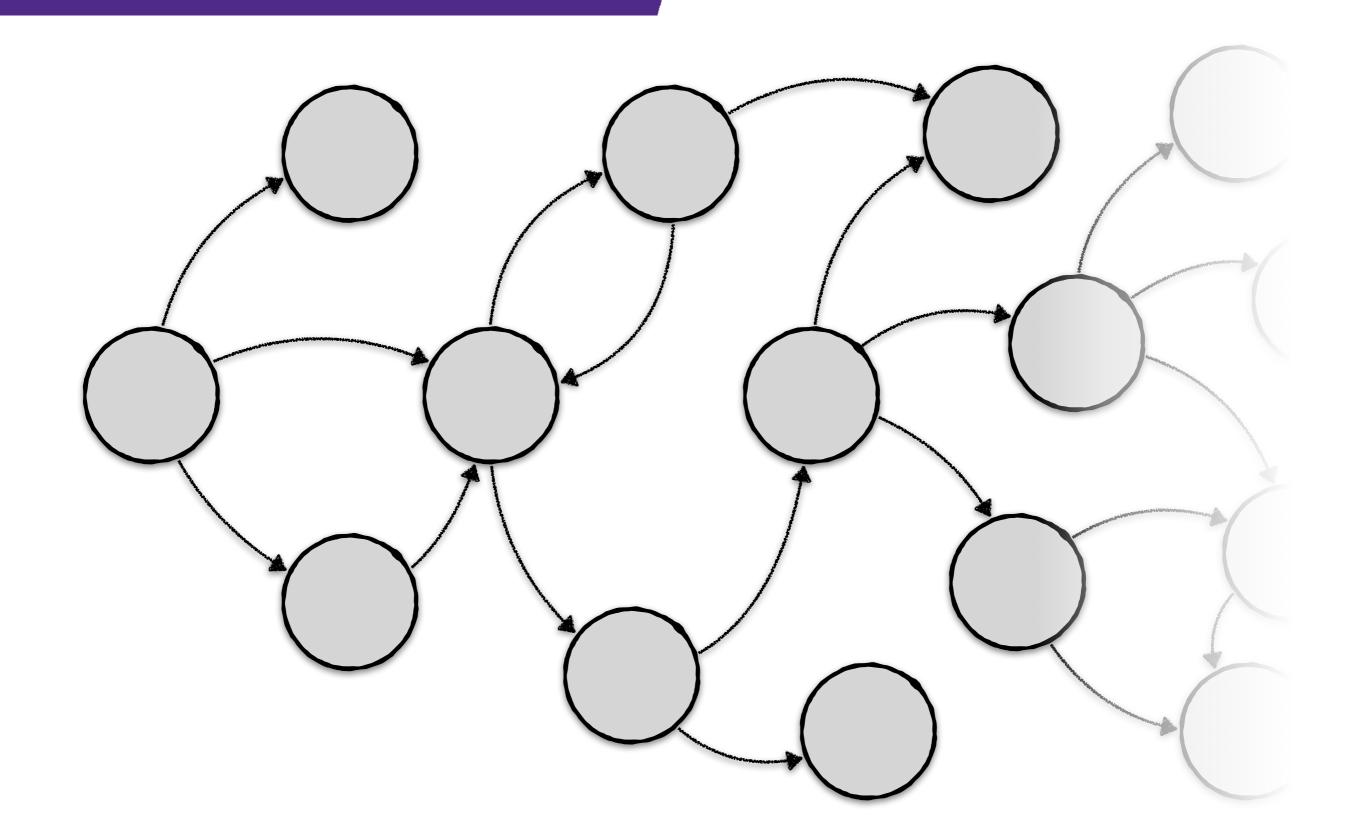
```
interface Client {
void sendCommand(Command command);
boolean hasResult();
Result getResult();
```

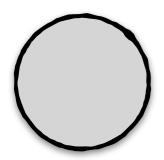
- A distributed system consists of a set of nodes which communicate over an asynchronous network, working together to run a protocol.
- Nodes are I/O automata; they run as single-threaded event loops.
- Nodes are split between client and server nodes.

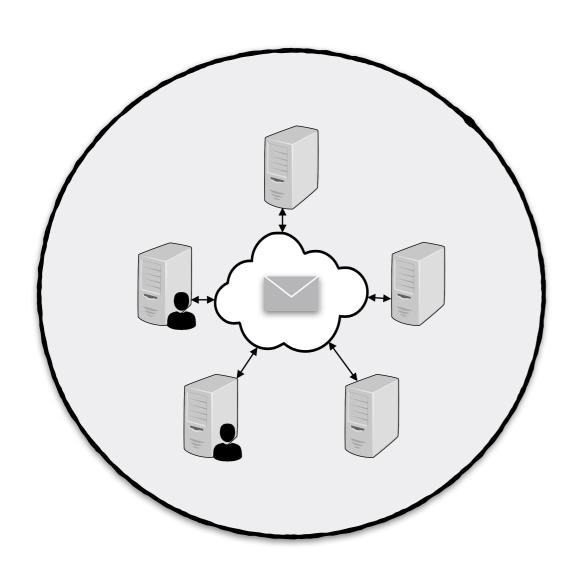


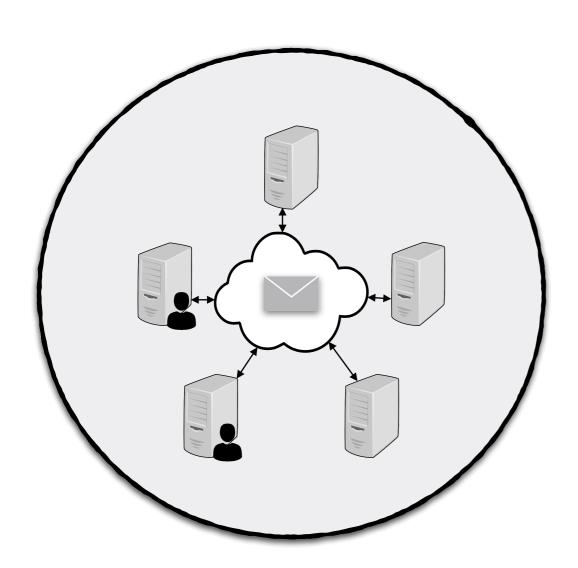
Programming Model Benefits

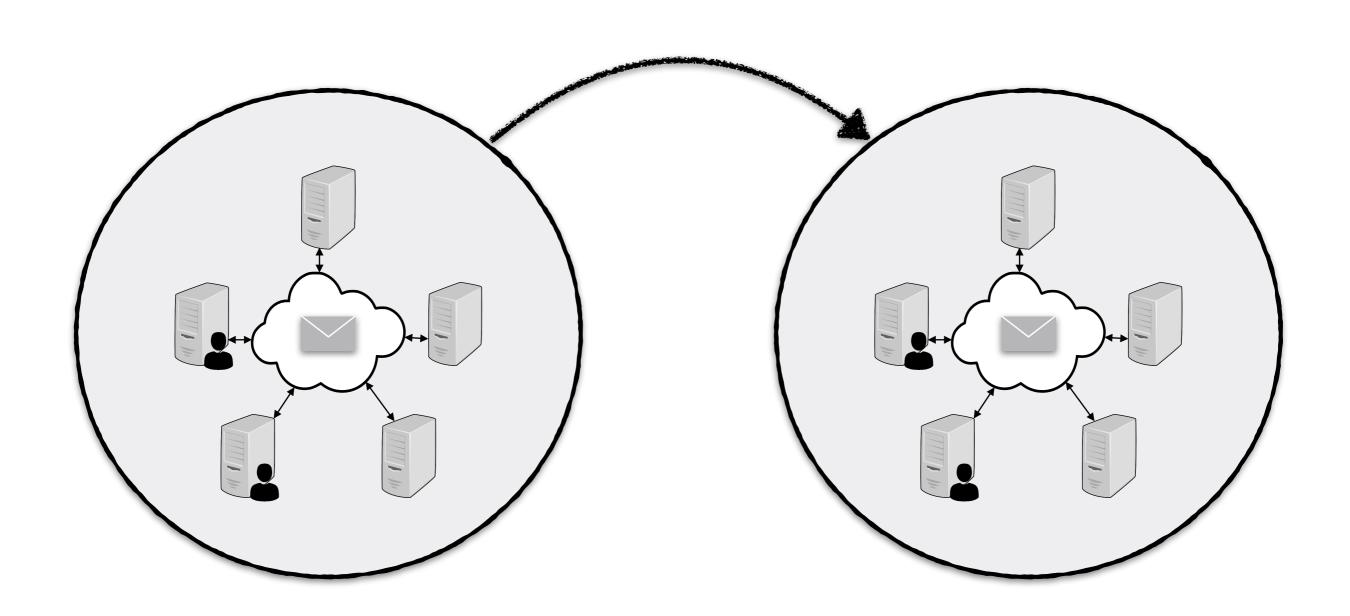
- Isolates concurrency to coarsest possible granularity
- Lets students focus on distributed protocols, avoiding issues such as deadlock within a node
- Allows for model checking at the protocol level without significant modification or overhead

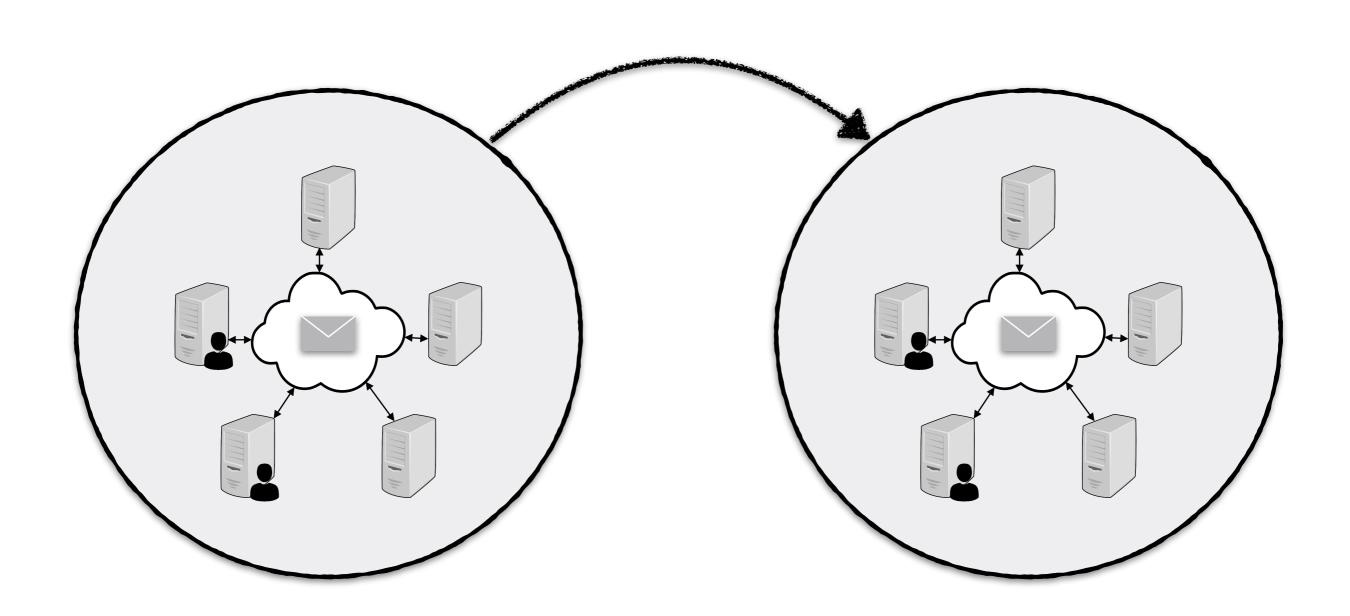


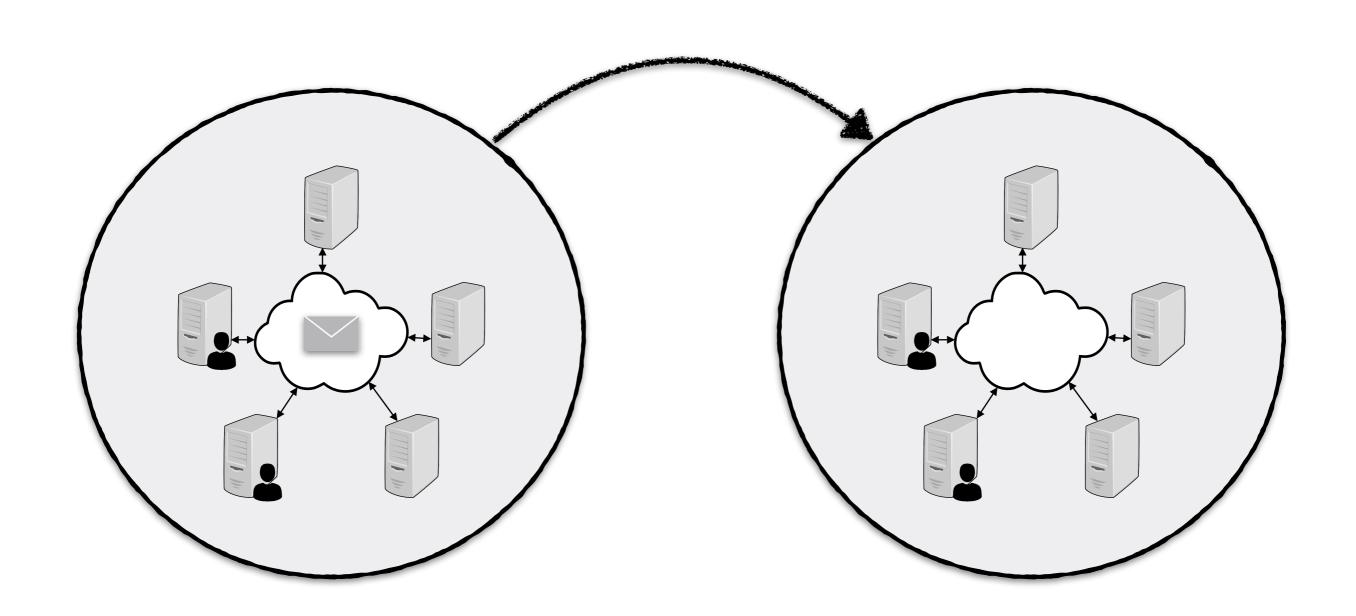


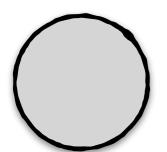


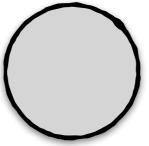


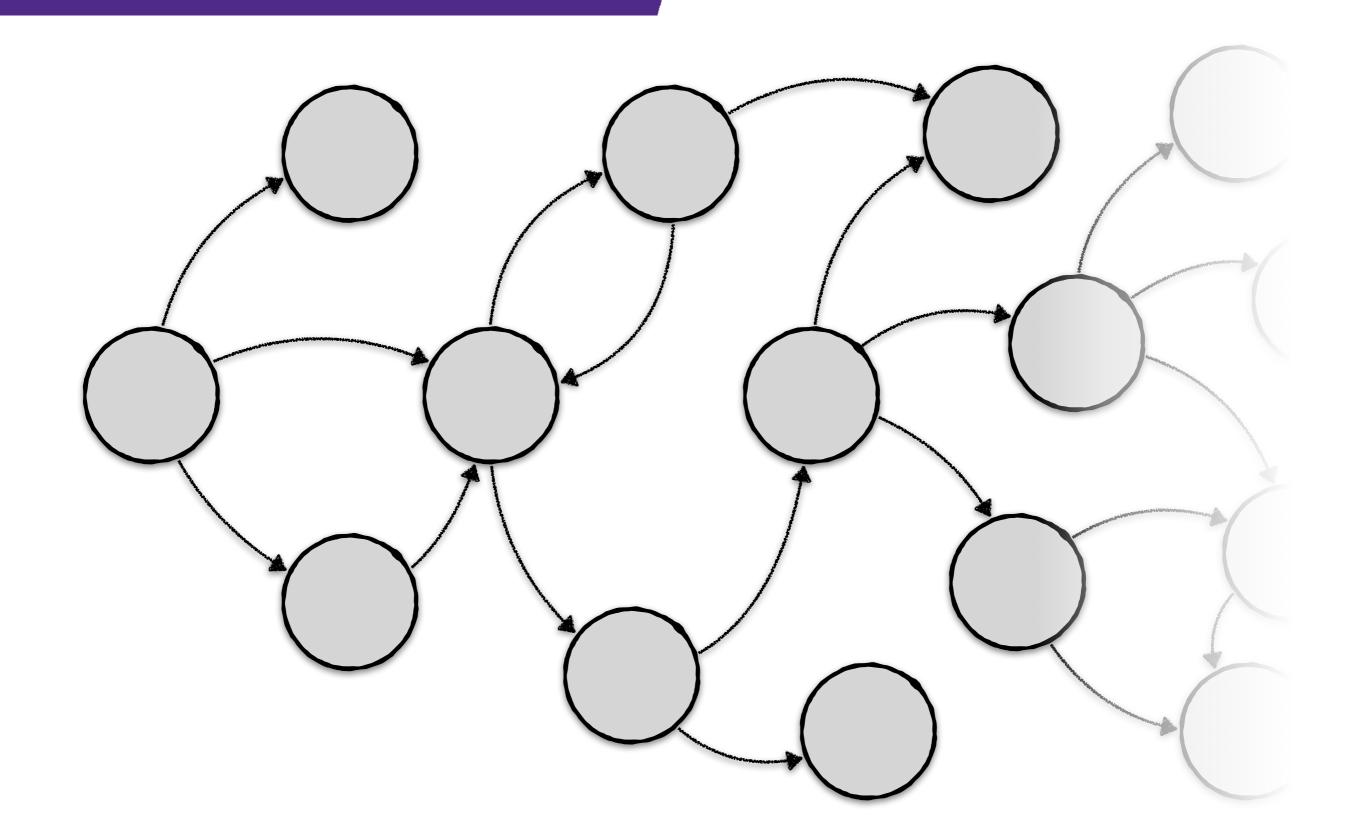


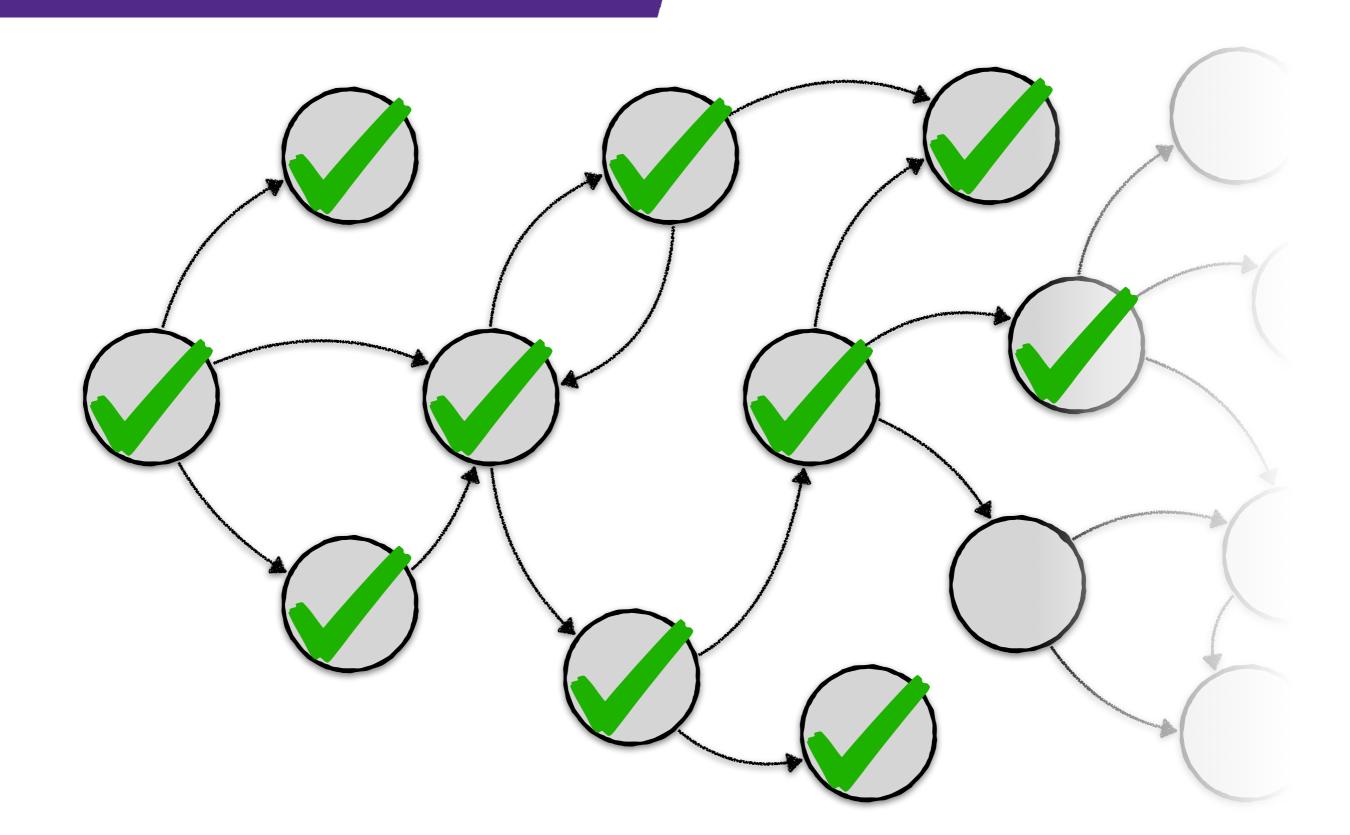


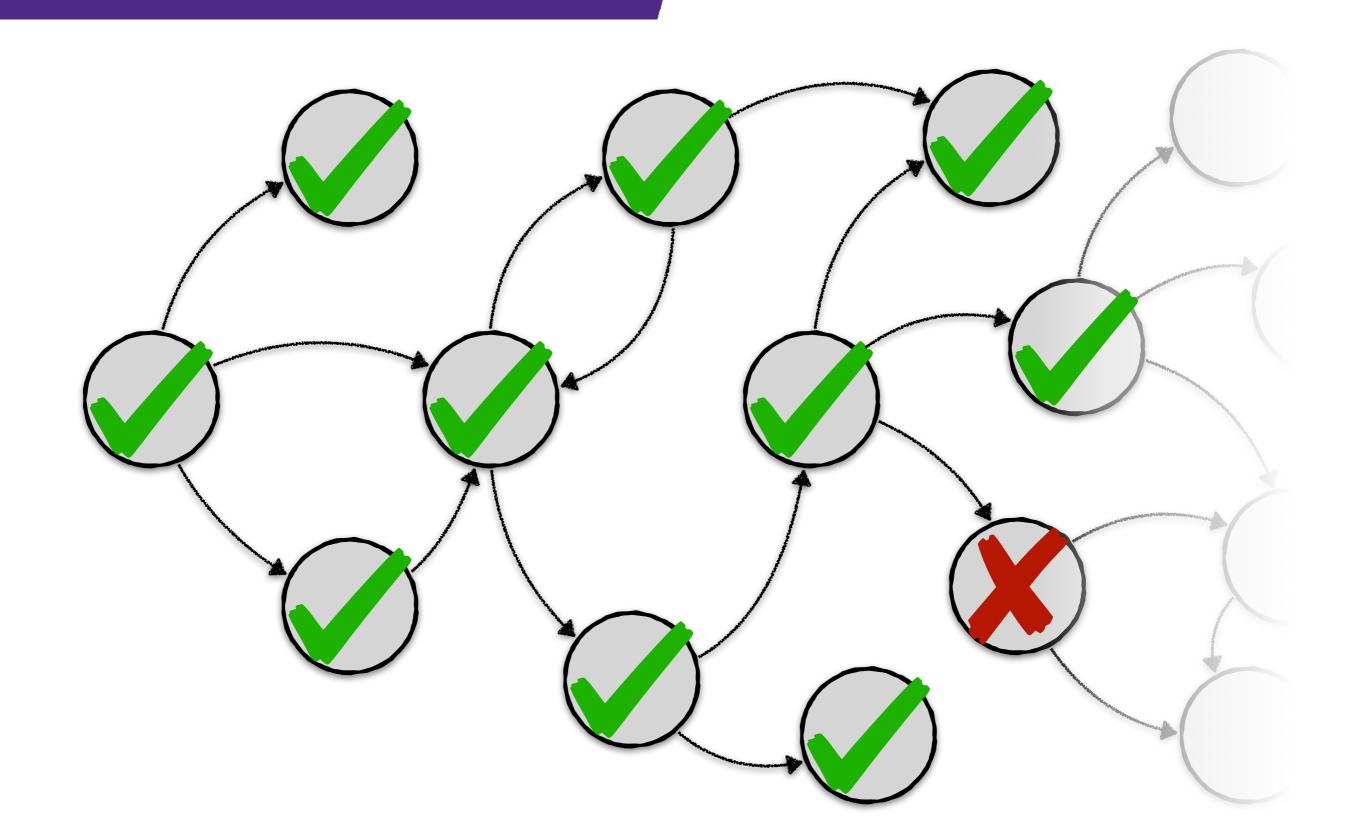












Outline

- 1. The DSLabs programming model
- 2. Model checking strategies and optimizations
- 3. Understandability and Oddity visual debugger
- 4. Experiences

How can the model checker evaluate states of student implementations?

What should the interface be between the tests and student implementations?

- Tests can check end-to-end properties, nothing else
- Allows maximum flexibility during implementation
- Doesn't allow checking more complicated properties, optimizations

- Tests can check end-toend properties, nothing else
- Allows maximum flexibility during implementation
- Doesn't allow checking more complicated properties, optimizations

- Message formats, and even internal data structures defined for students
- Allows for thorough, incremental checking
- Solves design challenges for students
- Couples tests to implementation

- Tests can check end-toend properties, nothing else
- Allows maximum flexibility during implementation
- Doesn't allow checking more complicated properties, optimizations

- Message formats, and even internal data structures defined for students
- Allows for thorough, incremental checking
- Solves design challenges for students
- Couples tests to implementation

- Tests can check endto-end properties, nothing else
- Allows maximum flexibility during implementation
- Doesn't allow checking more complicated properties, optimizations

Gray-Box

- Students implement limited, informational interface
- Allows enough insight into state for thorough checking
- Leaves most design decisions to students

- Message formats, and even internal data structures defined for students
- Allows for thorough, incremental checking
- Solves design challenges for students
- Couples tests to implementation

- Tests can check endto-end properties, nothing else
- Allows maximum flexibility during implementation
- Doesn't allow checking more complicated properties, optimizations

Gray-Box

- Students implement limited, informational interface
- Allows enough insight into state for thorough checking
- Leaves most design decisions to students

- Message formats, and even internal data structures defined for students
- Allows for thorough, incremental checking
- Solves design challenges for students
- Couples tests to implementation

Improving Model Checking Performance, Reliability

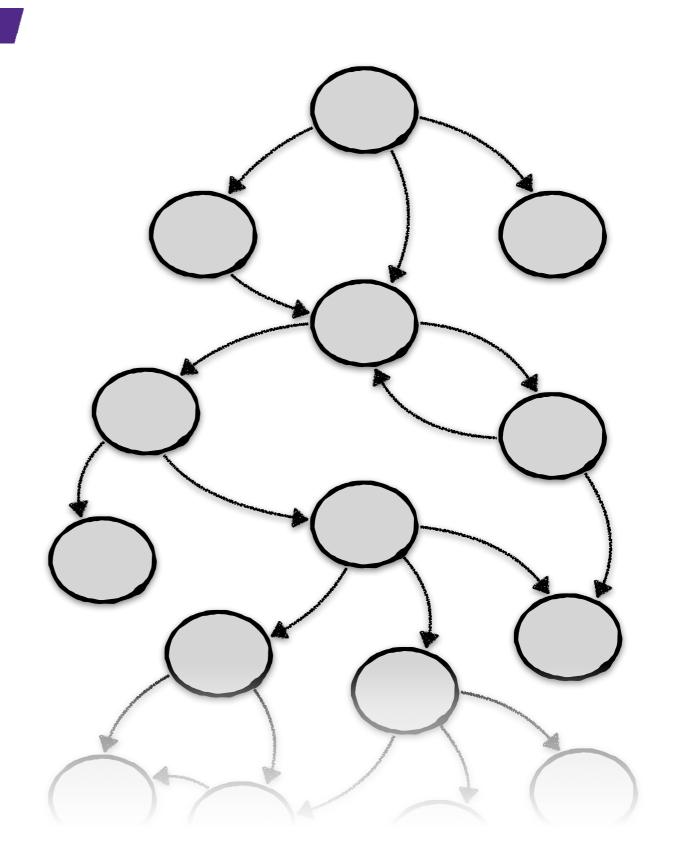
Model checking faces state-space explosion problem.

Strategies:

- 1. Pruning the search space
- 2. Punctuated search
- 3. Searching for progress

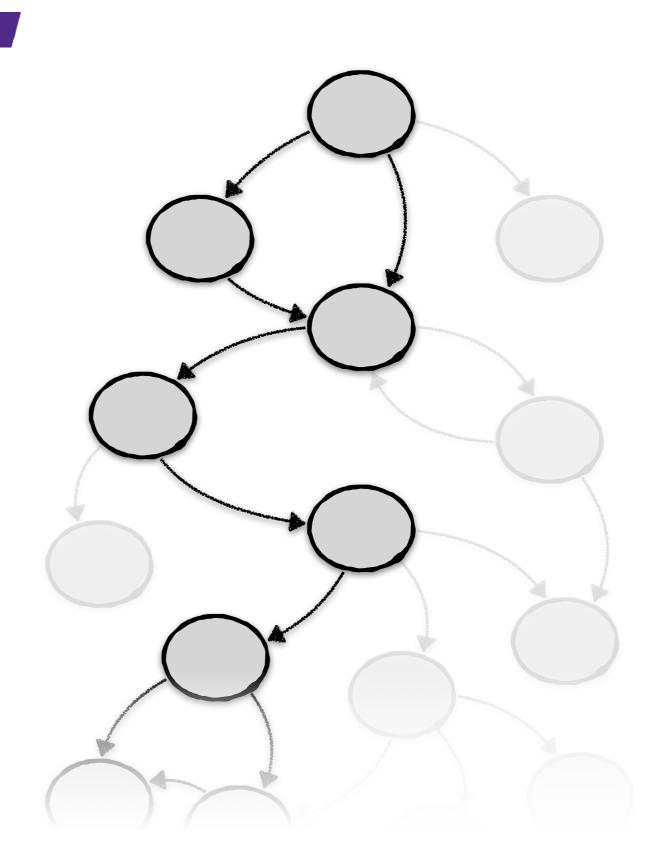
Pruning the Search Space

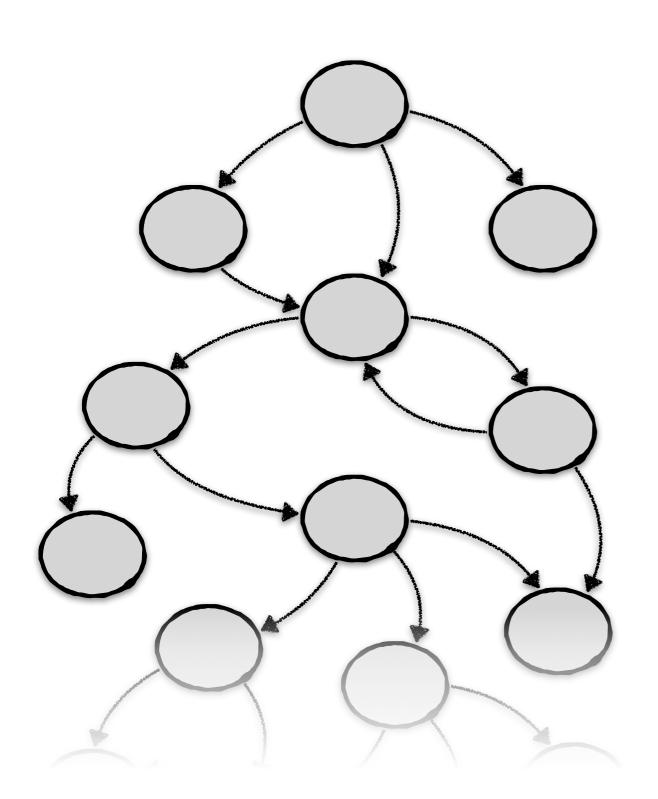
- Not all states are interesting.
- We can prune uninteresting states, refusing to expand them during the search.
- If we're interested in linearizability, we can safely ignore states in which clients have received all results.



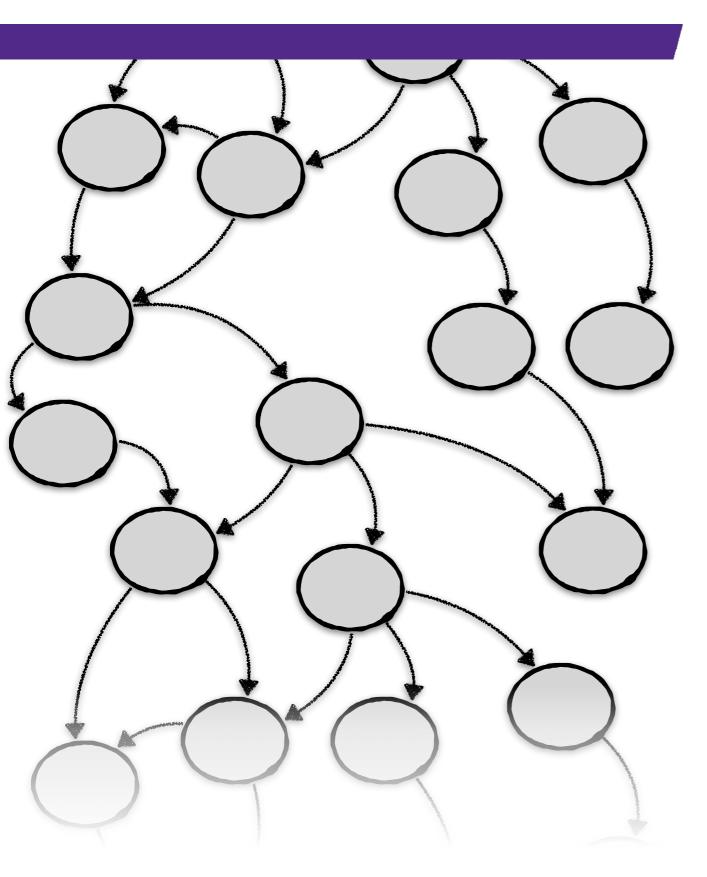
Pruning the Search Space

- Not all states are interesting.
- We can prune uninteresting states, refusing to expand them during the search.
- If we're interested in linearizability, we can safely ignore states in which clients have received all results.

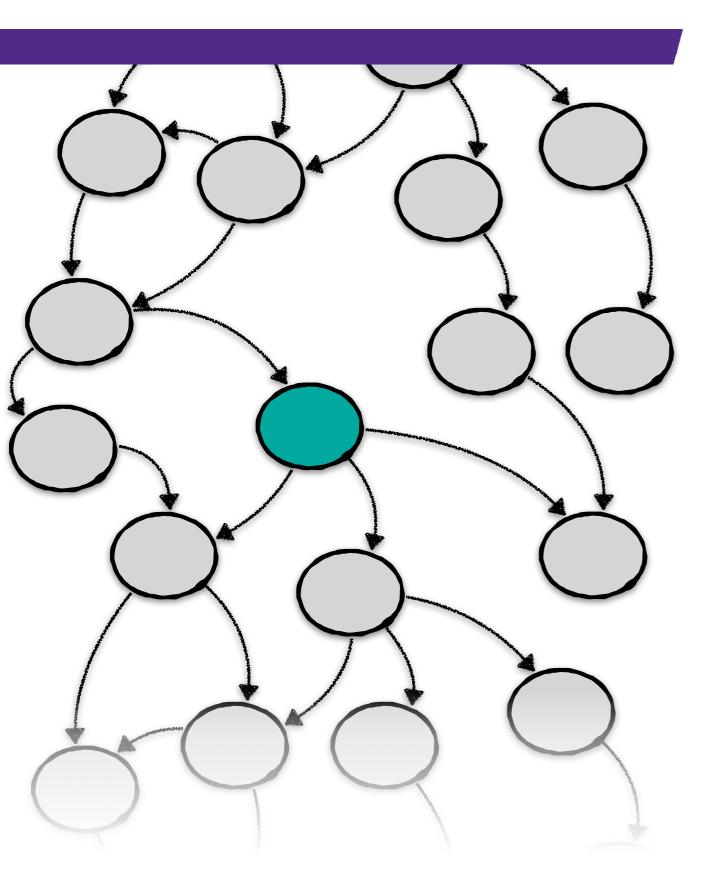




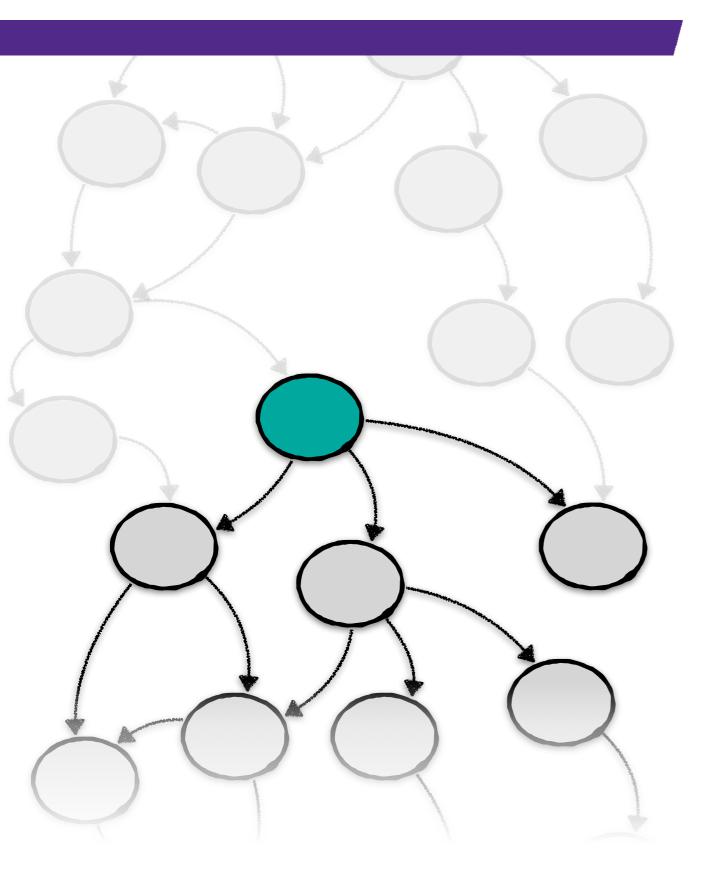
- BFS is limited primarily by the depth to which it can search.
- First, the model checker finds a state matching an intermediate constraint. Then, resumes checking starting from the new state.
- Repeatable, allows for scripting complex searches



- BFS is limited primarily by the depth to which it can search.
- First, the model checker finds a state matching an intermediate constraint. Then, resumes checking starting from the new state.
- Repeatable, allows for scripting complex searches

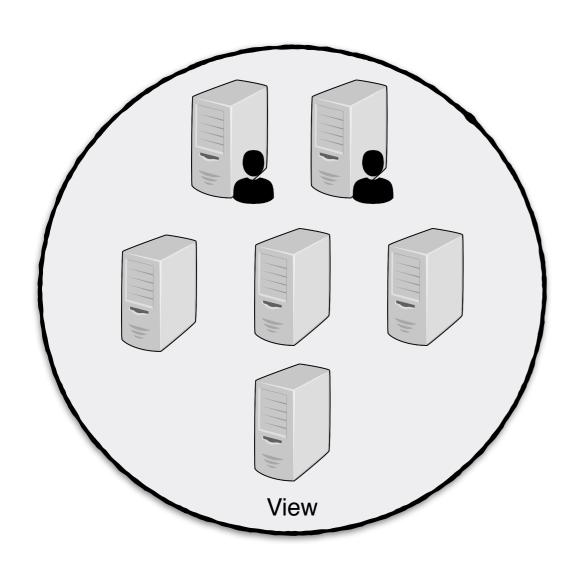


- BFS is limited primarily by the depth to which it can search.
- First, the model checker finds a state matching an intermediate constraint. Then, resumes checking starting from the new state.
- Repeatable, allows for scripting complex searches

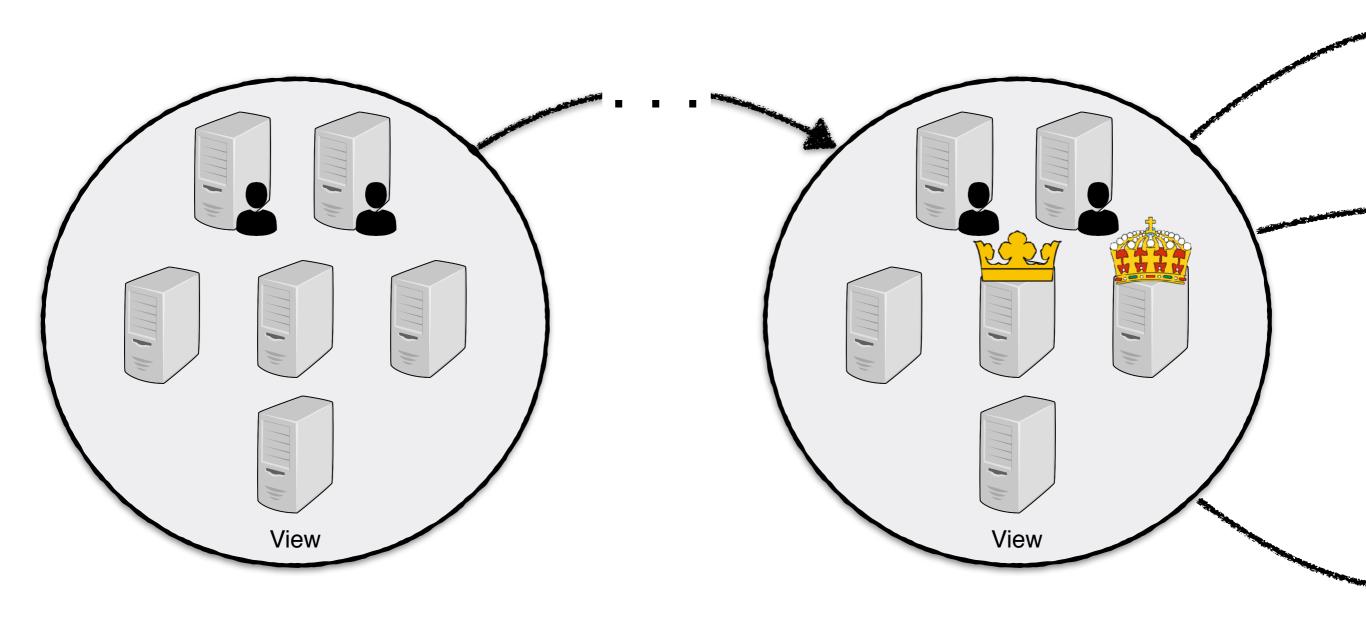


- BFS is limited primarily by the depth to which it can search.
- First, the model checker finds a state matching an intermediate constraint. Then, resumes checking starting from the new state.
- Repeatable, allows for scripting complex searches

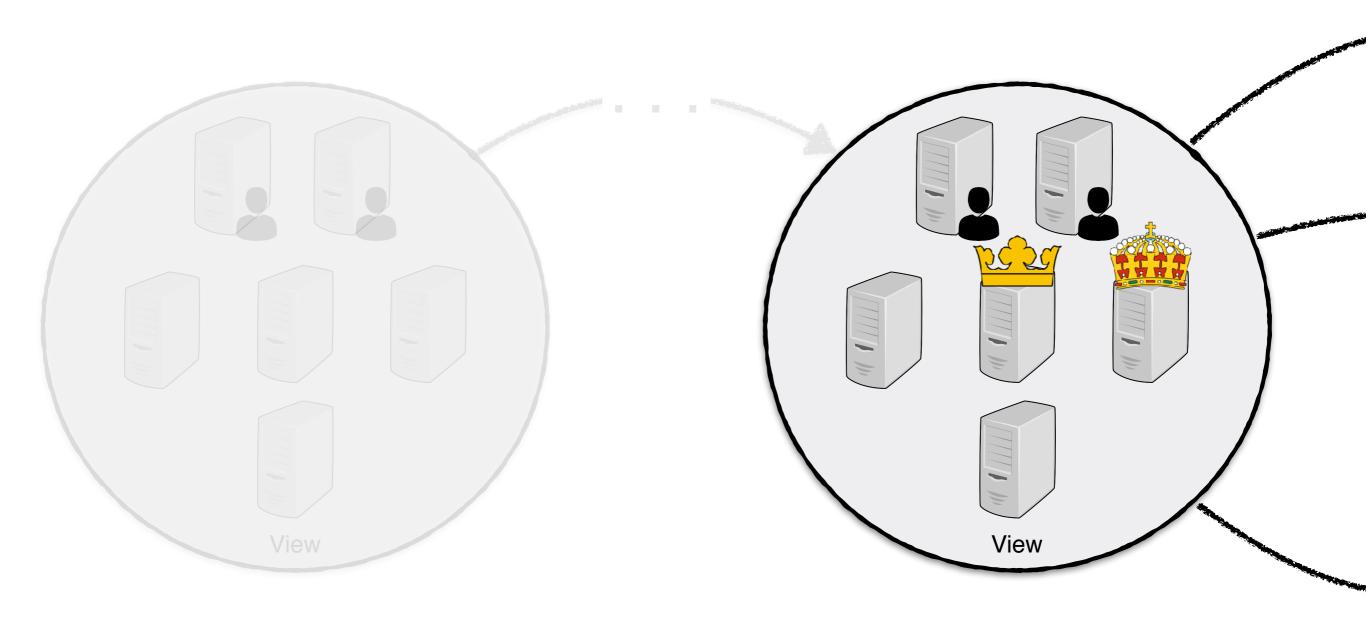
Punctuated Search Example: Primary-backup



Punctuated Search Example: Primary-backup



Punctuated Search Example: Primary-backup



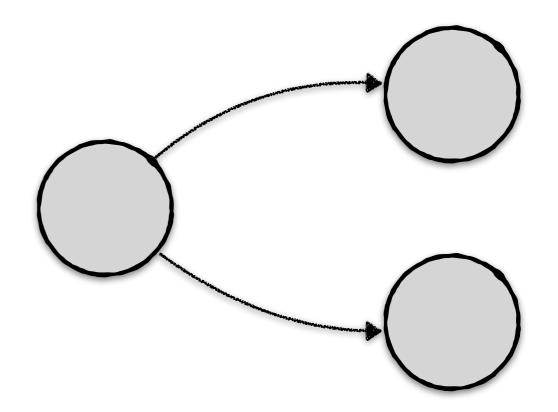
Simplifying Implementation: Testing Determinism

- Key assumption: nodes are deterministic.
- Some sources of nondeterminism are non-obvious.
- DSLabs has flag to check handler determinism, facilitating correct implementation.



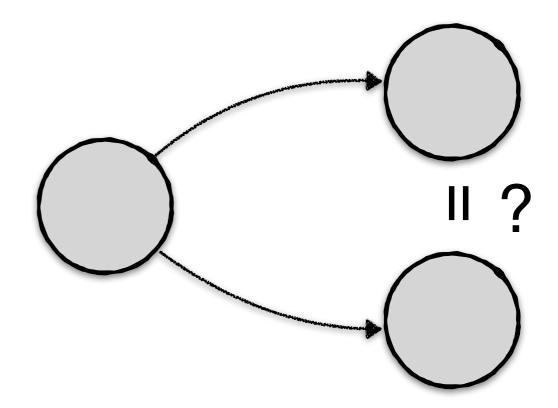
Simplifying Implementation: Testing Determinism

- Key assumption: nodes are deterministic.
- Some sources of nondeterminism are non-obvious.
- DSLabs has flag to check handler determinism, facilitating correct implementation.



Simplifying Implementation: Testing Determinism

- Key assumption: nodes are deterministic.
- Some sources of nondeterminism are non-obvious.
- DSLabs has flag to check handler determinism, facilitating correct implementation.



Designing Systems for Model Checking

- Performance of model checking is implementationdependent; runtime optimizations can reduce checkability.
- Our advice to students:
 - Favor simplicity.
 - Keep and send minimal state.
 - Ensure system can make progress with minimal steps.

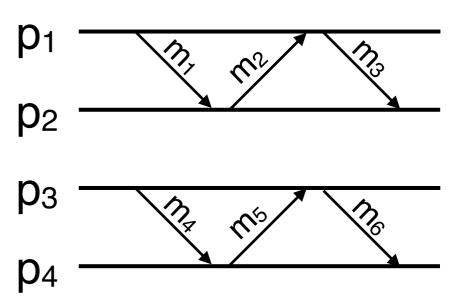
Outline

- 1. The DSLabs programming model
- 2. Model checking strategies and optimizations
- 3. Understandability and Oddity visual debugger
- 4. Experiences

- A trace is a linearization of an execution returned by model checker, demonstrating invariant violation.
- BFS used by model checker could return any minimal length trace.
- DSLabs performs a depth-first topological sort of the event graph before returning traces to students



- A trace is a linearization of an execution returned by model checker, demonstrating invariant violation.
- BFS used by model checker could return any minimal length trace.
- DSLabs performs a depth-first topological sort of the event graph before returning traces to students



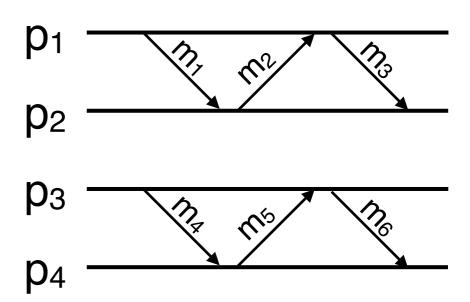
1. m_1 m_2 m_3 m_4 m_5 m_6

- A trace is a linearization of an execution returned by model checker, demonstrating invariant violation.
- BFS used by model checker could return any minimal length trace.
- DSLabs performs a depth-first topological sort of the event graph before returning traces to students



1. m_1 m_2 m_3 m_4 m_5 m_6

- A trace is a linearization of an execution returned by model checker, demonstrating invariant violation.
- BFS used by model checker could return any minimal length trace.
- DSLabs performs a depth-first topological sort of the event graph before returning traces to students

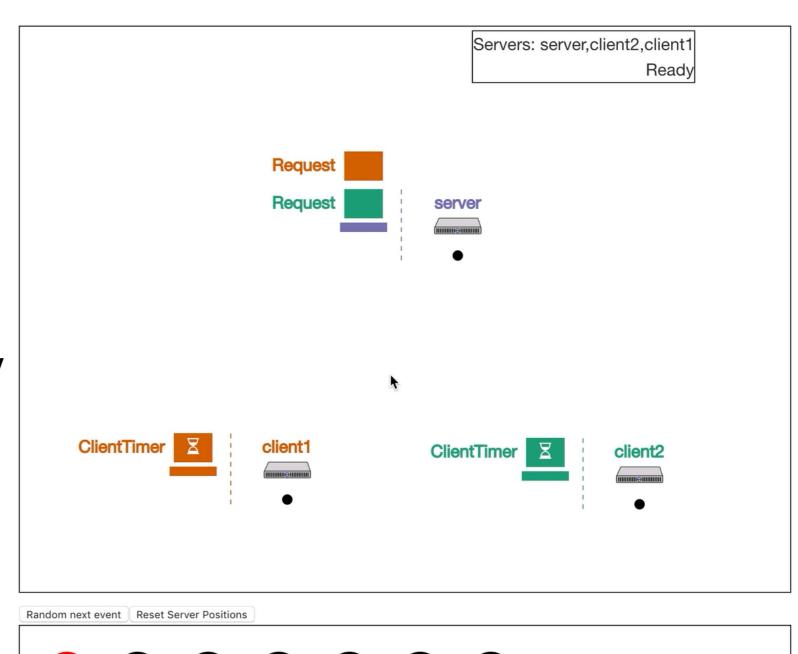


1. m_1 m_2 m_3 m_4 m_5 m_6

2. m_1 m_4 m_5 m_2 m_6 m_3

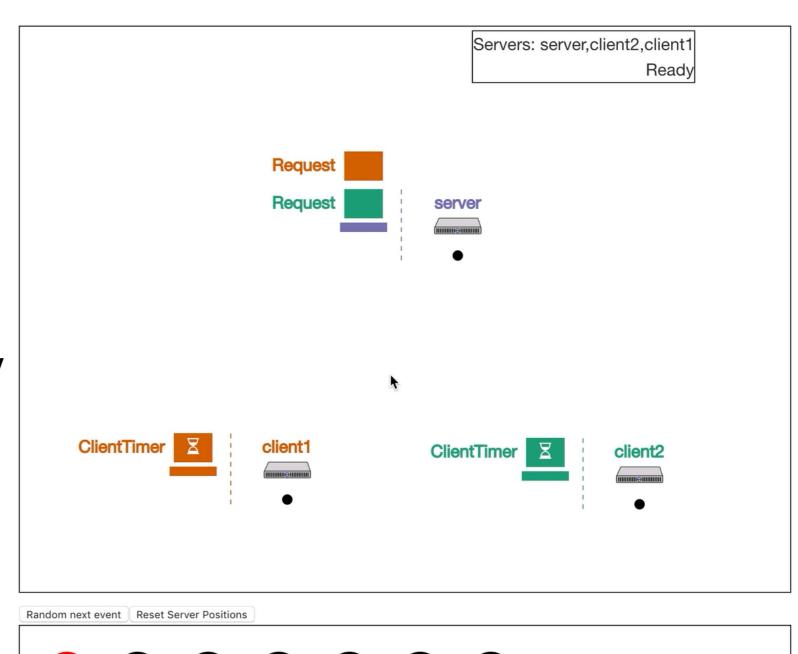
Oddity

- Allows exploration from initial state or invariantviolating trace
- Lets students interactively explore states, examine messages and nodes
- Can "time-travel," explore alternate histories



Oddity

- Allows exploration from initial state or invariantviolating trace
- Lets students interactively explore states, examine messages and nodes
- Can "time-travel," explore alternate histories

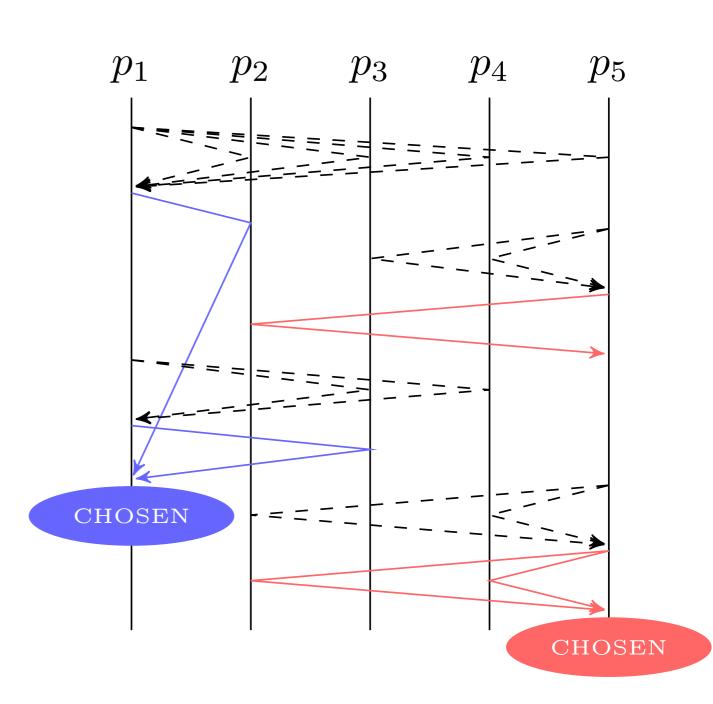


Outline

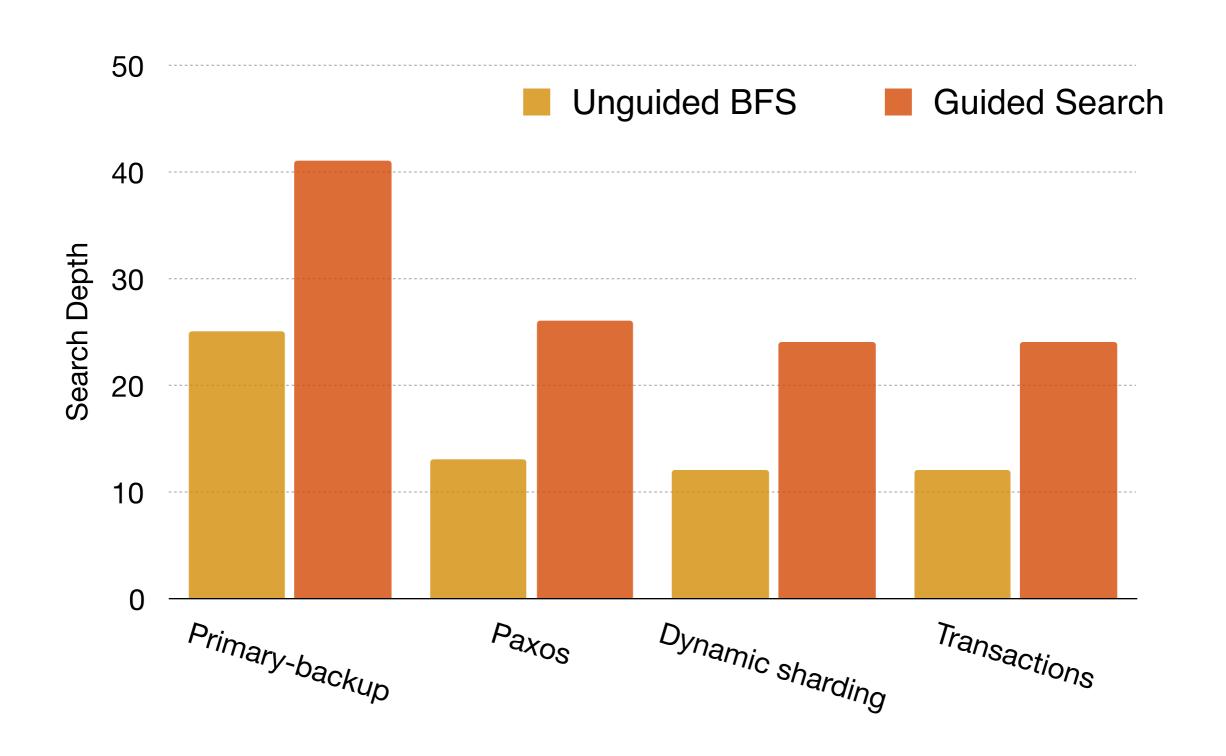
- 1. The DSLabs programming model
- 2. Model checking strategies and optimizations
- 3. Understandability and Oddity visual debugger
- 4. Experiences

Can Guided Searches Find Bugs?

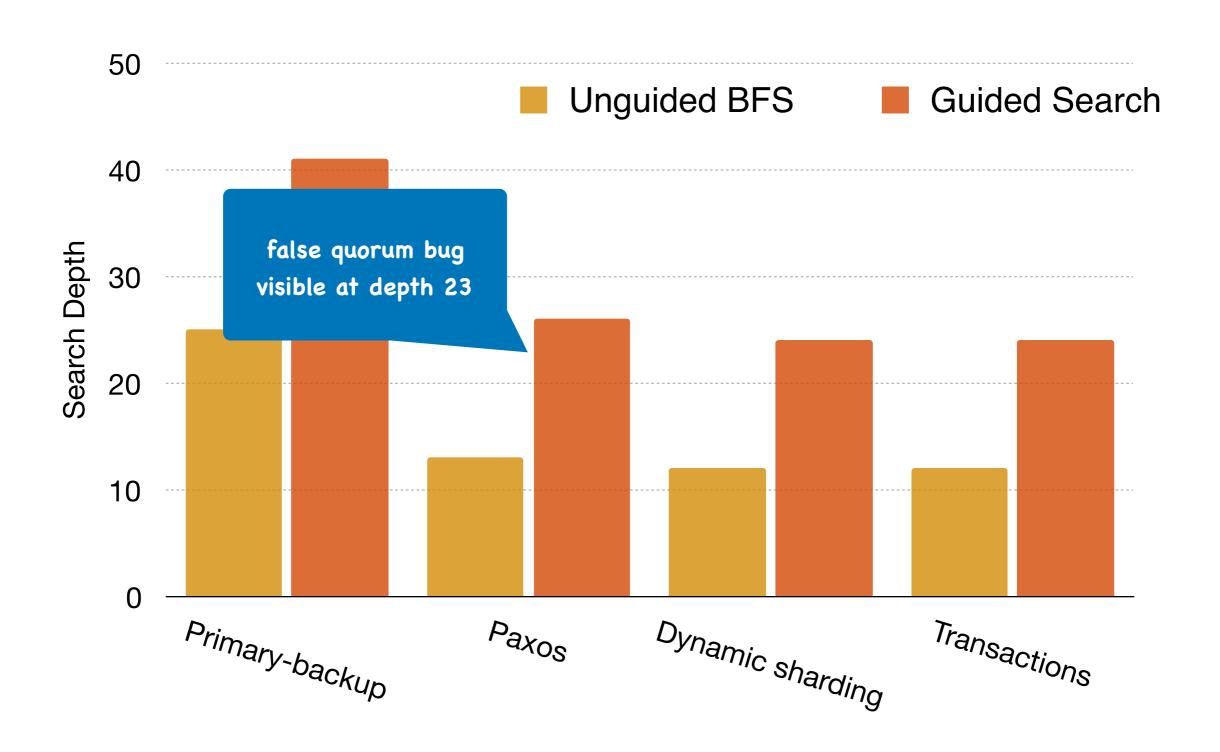
- Naïve BFS can't find the example false quorum bug.
- Random exploration takes an average of 12 hours.
- Guided search for this type of bug takes just 18 seconds.



Can Guided Search Improve Model Checking Thoroughness?



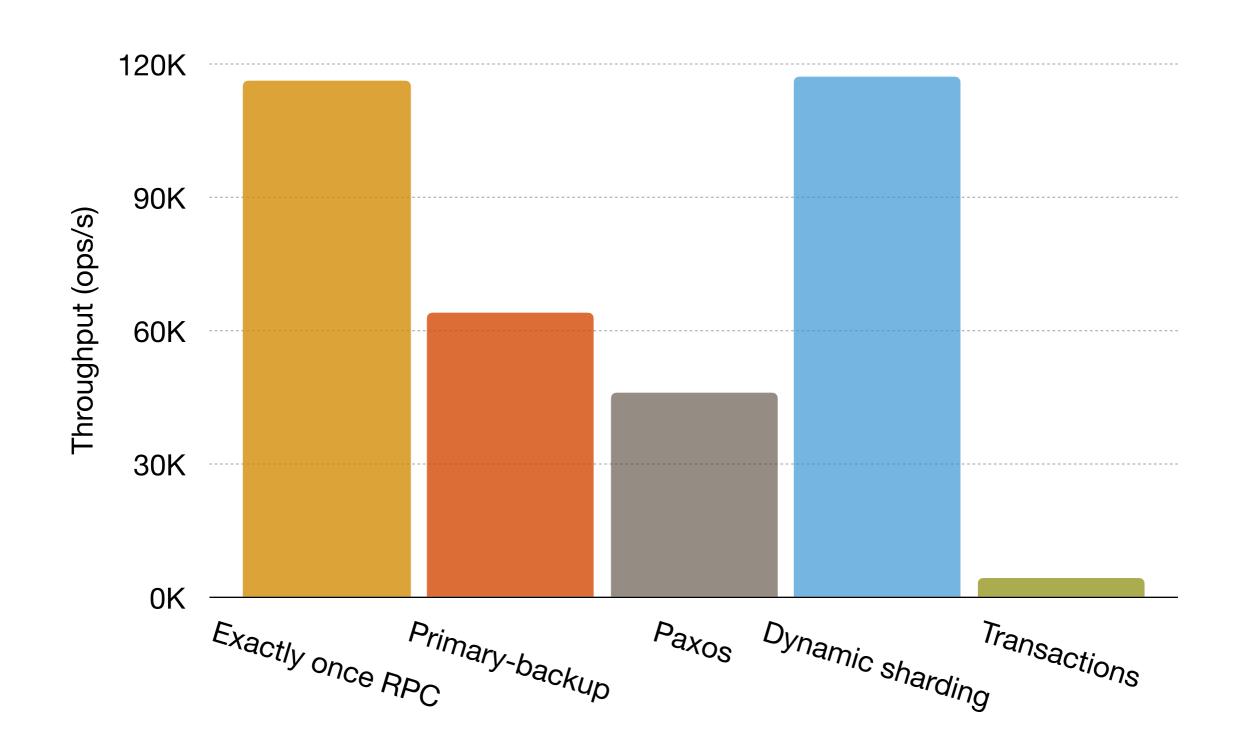
Can Guided Search Improve Model Checking Thoroughness?



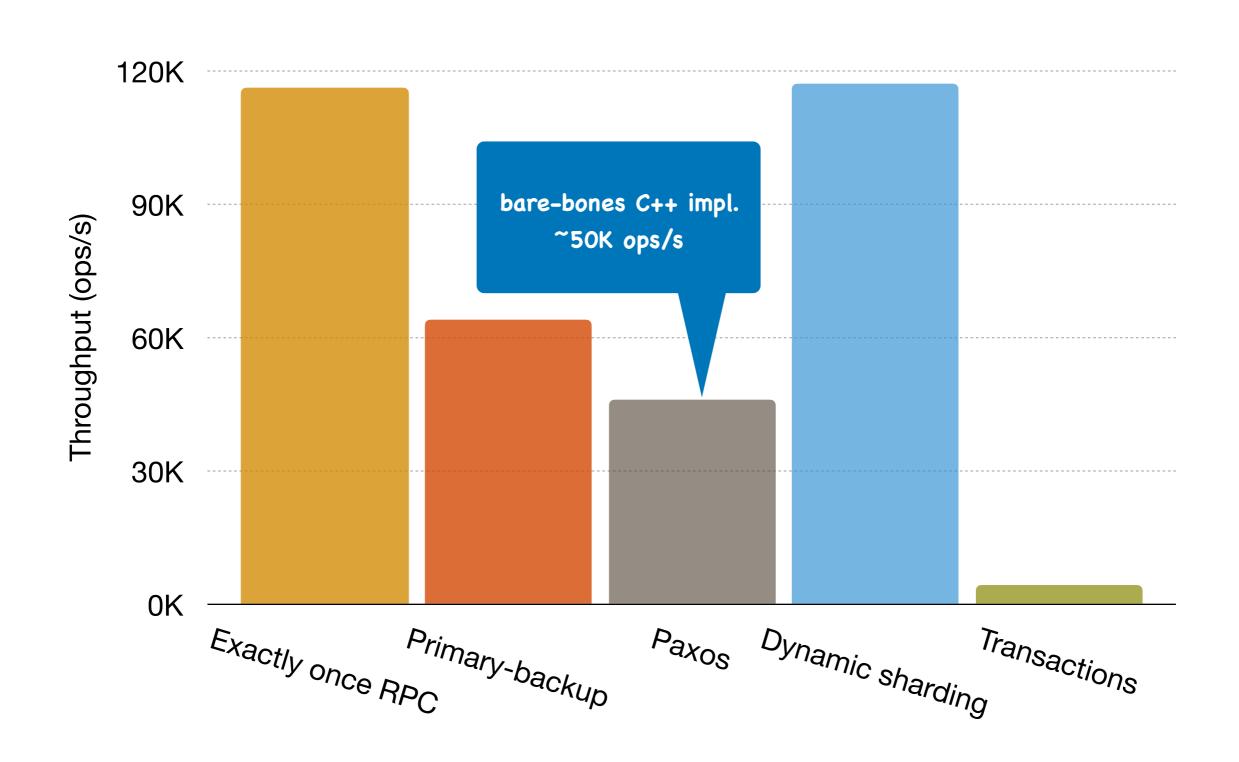
Are Students Able to Debug Their Systems?

- Based on opt-in telemetry: over 150 invariant-violations examined with Oddity
- Almost all of these fixed before submission
- Only 25 submissions (across all assignments) found to violate invariants, 38 unable to pass searches for progress

Can Students Build Runnable, Performant Systems?



Can Students Build Runnable, Performant Systems?



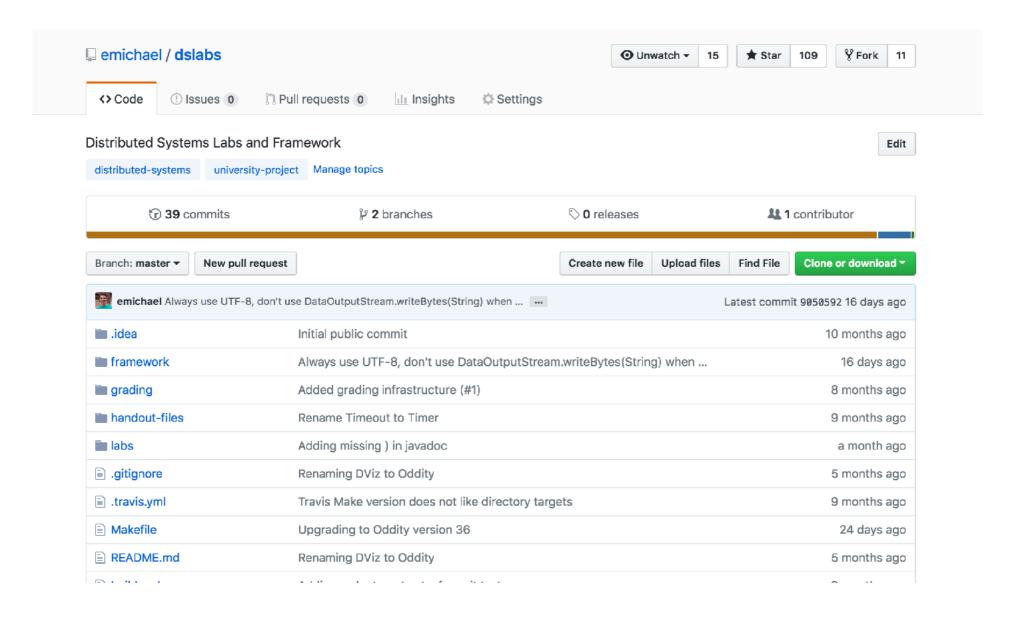
Does DSLabs Encourage "Distributed Thinking"?

- We want to encourage a distributed systems mindset: focus on invariants, rather than normal case.
- Model checking centers the distributed programming environment, finds "rare" errors.
- Visual debugger reinforces the programming model.

Summary

- DSLabs, a new framework for building distributed systems assignments:
 - Uses efficient model checking based on guided search techniques,
 - Allows instructors to design model checking tests for student implementations,
 - Includes tools for debugging, understanding errors when they occur.
- DSLabs has been invaluable at UW, helped us scale undergraduate distributed systems to 200 students per quarter.

Thanks for Listening!





https://github.com/emichael/dslabs

emichael@cs.washington.edu

Feedback, issues, pull-requests welcome